

Linear Programming in Practice

▪ **Essential Issue:** To model non-linear reality with linear equations

- Activities
- Piece-wise linear approximations
- Fixed charges

▪ **Another practical question**

- Duality

Activities

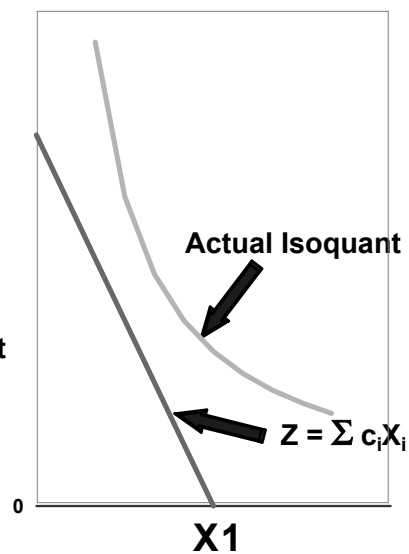
▪ **Motivation:**

▪ **If we use a standard production function**

$$f(\underline{X}) = \sum c_i X_i = Z$$

resources, \Rightarrow output

▪ **We are not able to represent typical production function with diminishing marginal returns and non-linear isoquants**



Activities (continued)

▪Concept

- An activity is a
 - Specific way to use resources
 - in fixed proportions

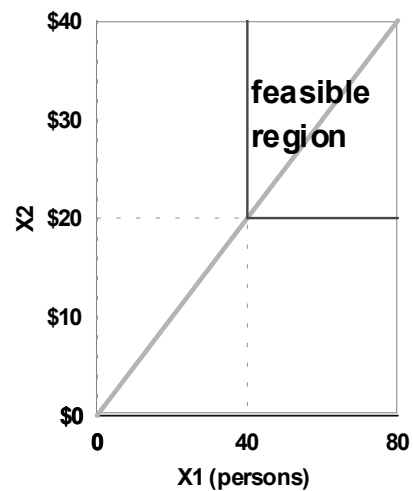
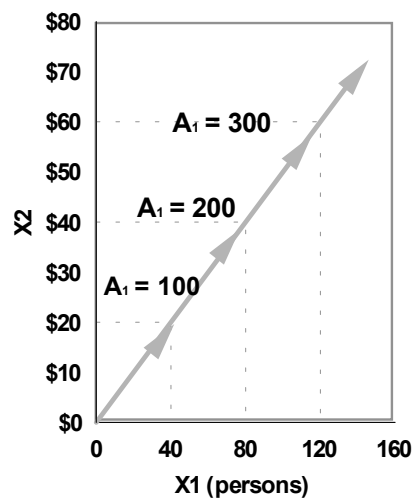
▪Physical interpretation is direct, e.g.:

- an aircraft using pilots, fuel / ton-km
- a machine requiring labor, materials per unit product

▪Think of activities as intermediates between resources and output

resources  activities  output

Example: transport process A_1 uses 40 persons, \$20k to produce 100 T-km



Activities

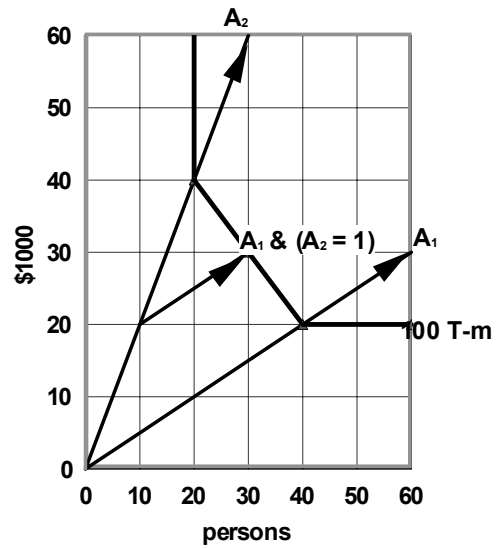
▪ Two Activities

▪ $A_1 = (40, 20K)$
 $\implies 100 \text{ T-m}$

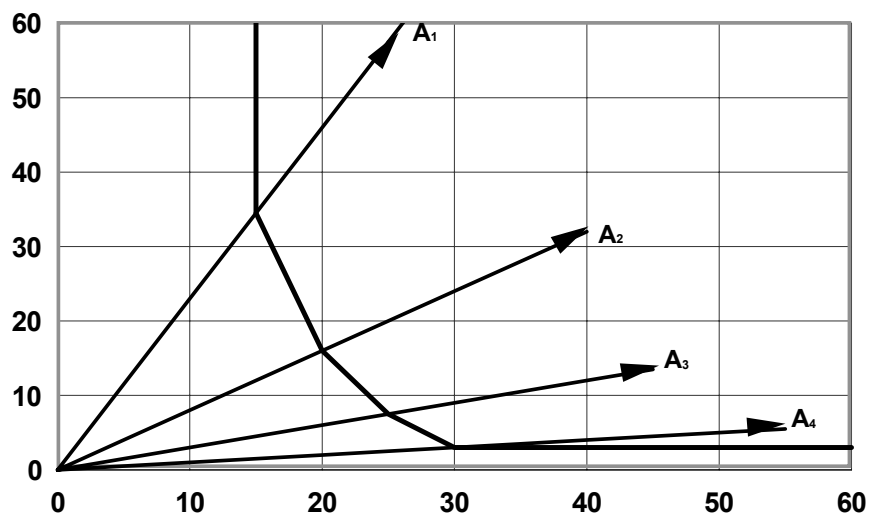
▪ $A_2 = (10, 20K)$
 $\implies 50 \text{ T-m}$

$A_2 = 1 \{10, 20\} \implies 50$

$A_1 = \frac{1}{2} \{20, 10\} \implies 50$



Many Activities



LP Formulation with Activities

	Cr (kg)	C (kg)	Profit (\$)
Process 1	6	4	30
Process 2	5	2	28
Process 3	3	6	29

▪Optimize: $Z = \sum c_i A_i$ -- subject to constraints

▪Example: Alloy optimization

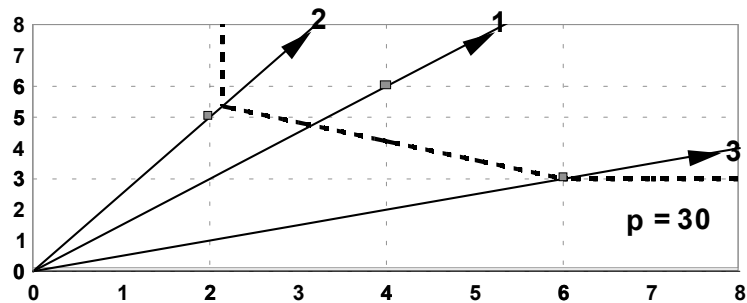
–Three possible processes, each with different unit costs

–subject to limits on Cr, C content

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Formulation



$$\max Z = 30 P_1 + 28 P_2 + 29 P_3$$

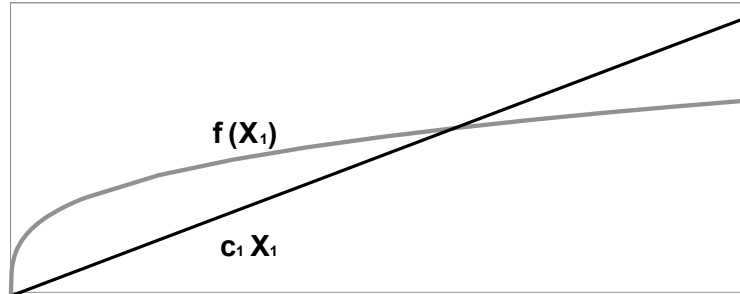
$$\text{s.t. } 6 P_1 + 5 P_2 + 3 P_3 \leq 26 \text{ (Cr)}$$

$$4 P_1 + 2 P_2 + 6 P_3 \leq 7 \text{ (C)}$$

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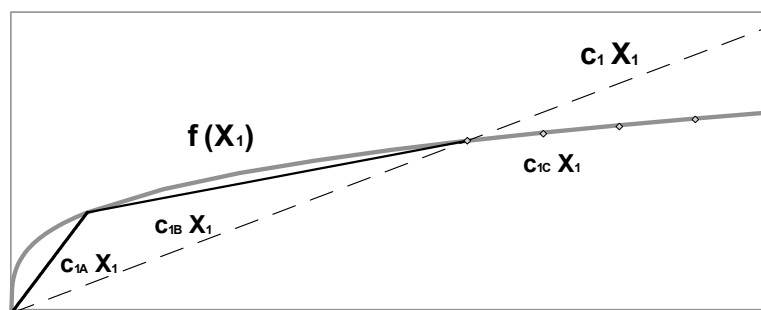
Piece-Wise Linear Approximations (1)



▪Motivation:

- Returns to scale generally non-linear
- Straight line approximations are inaccurate

Piece-Wise Linear Approximations (2)



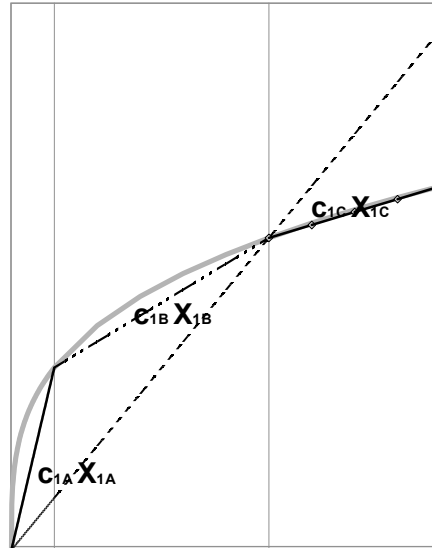
▪Concept:

- Represent $f(X_1)$ with several lines

Piece-wise Linear Approximations

Implementation Notes:

- X_1 must be redefined as several variables - X_{1A}, X_{1B}, \dots
- These new variables must not overlap, so $X_{1A} < X_{1B}$, etc.
- New variables and constraints make the LP larger and, therefore, more expensive



Piece-wise Linear Approximations

Mathematically:

▪ Given:
$$\begin{aligned} \text{Max } Z &= f(X_1) + 4X_2 \\ \text{s.t.} \quad & 3X_1 + 6X_2 \leq 8 \end{aligned}$$

Piece-wise linear approximation gives:

- $X_1 \Rightarrow X_{1A} + X_{1B}$
- X_{1A}, X_{1B} have same a_{ij} as X_1
- $c_1 = c_{1A}, c_{1B}$
- $X_{1A} < \text{cutoff } X \text{ value between } X_{1A} \text{ and } X_{1B}, X'$

▪ Thus:
$$\begin{aligned} \text{Max } Z &= c_{1A} X_{1A} + c_{1B} X_{1B} + 4X_2 \\ \text{s.t.} \quad & 3 X_{1A} + 3X_{1B} + 6X_2 \leq 8 \\ & X_{1A} \leq X' \end{aligned}$$

Piece-wise Linear Approximations (4)

- **Key Limitation:**

- ONLY works for convex feasible region!

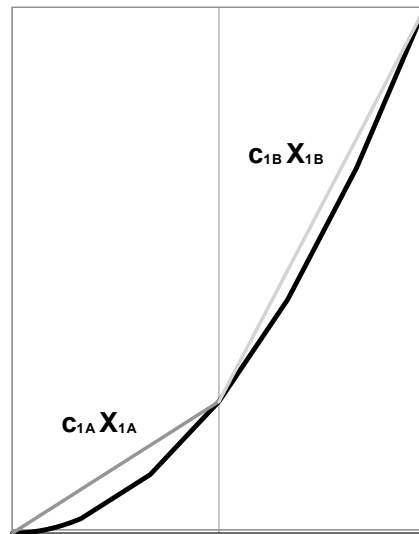
- **Why?**

- What if $c_{1B} > c_{1A}$? (see fig)

$$\text{Max } Z = c_{1A} X_{1A} + c_{1B} X_{1B} + 4X_2$$

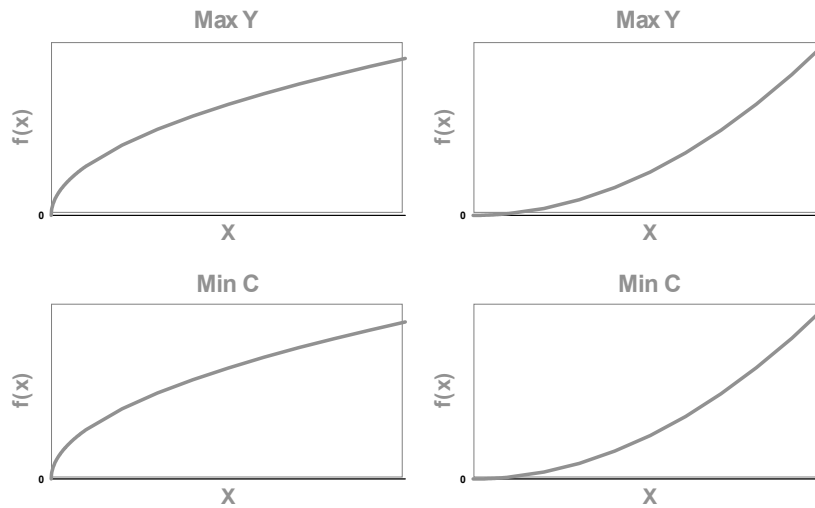
- **The LP will select X_{1B} before X_{1A} ;**

Result may be meaningless!



Convex Feasible Regions Review:

Piecewise linear approximation works when FR is convex



Fixed Charges

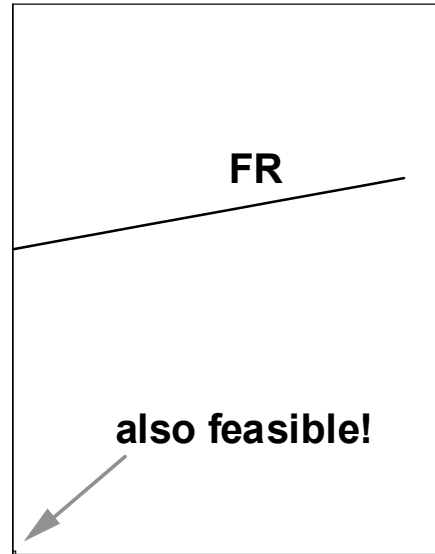
- Example: Warehousing
 - Cost = fixed rent, etc. + variable
 - Unless you choose not to operate it!

$$\begin{aligned} f(X_1) &= c_0 + c_1 X_1 & X_1 &\geq 0 \\ f(X_1) &= 0 & X_1 &= 0 \end{aligned}$$

- LP generally cannot handle fixed charges

Exception:

- All $X_i > 0$; $X_i \neq 0$
- then subtract $\sum c_0$ and optimize



Duality

▪ Concept:

- A "dual" is a mirror-image form to another problem (the "primal")
- If primal = max; then dual = min
- If primal = min; then dual = max
- Dual contains all information of the primal, but in a different format
- Optimum value of primal = optimum value of dual

▪ Example:

- Primal: maximize output subject to budget limitations
- Dual: minimize costs subject to output requirements

LP Duality

▪Mathematics:

–Given a Primal:

$$\begin{array}{ll} \text{Optimize:} & Z = \underline{c} X \\ \text{subject to:} & \underline{A} X \leq \geq \underline{B} \end{array}$$

–Dual is:

$$\begin{array}{ll} \text{Optimize:} & Y = \underline{B}^T W \\ \text{subject to:} & \underline{A}^T W \leq \geq \underline{c}^T \end{array}$$

▪Change of dimensionality between primal & dual:

– \underline{c}^T and \underline{B} have different number of variables

▪Can use duality to:

- Reduce size of constraint matrix
- Speed up LP solution

LP Duality - Example

$$\begin{array}{ll} \text{–Primal:} & \text{Max: } Z = X_1 + 2X_2 + 3X_3 \\ & \text{s.t. } 4X_1 + 2X_2 \leq 5 \\ & \quad 6X_1 + 7X_2 + 9X_3 \leq 12 \end{array}$$

$$\begin{array}{ll} \text{–A =} & \begin{array}{ccc} 4 & 2 & 0 \\ 6 & 7 & 9 \end{array} & \text{A}^T = \begin{array}{c} 4 \ 6 \\ 2 \ 7 \\ 0 \ 9 \end{array} \end{array}$$

$$\begin{array}{ll} \text{–B =} & \begin{array}{c} 5 \\ 12 \end{array} & \text{B}^T = \begin{array}{cc} 5 & 12 \end{array} \end{array}$$

$$\begin{array}{ll} \text{–C =} & \begin{array}{ccc} 1 & 2 & 3 \end{array} & \text{C}^T = \begin{array}{c} 1 \\ 2 \\ 3 \end{array} \end{array}$$

LP Duality - Example - continued

–Primal: Max: $Z = X_1 + 2X_2 + 3X_3$
 s.t. $4X_1 + 2X_2 \leq 5$
 $6X_1 + 7X_2 + 9X_3 \leq 12$

–Dual: Min: $Y = 5W_1 + 12W_2$
 s.t. $4W_1 + 6W_2 \geq 1$
 $2W_1 + 7W_2 \geq 2$
 $9W_2 \geq 3$

LP Duality - Interpretation of Results

–Primal:

Max: $Z = 3X_1 + X_2 + 8X_3$			
s.t.	$X_1 + X_3$	\leq	4
	$X_1 + X_2 + X_3$	\leq	7
	$2X_2 + X_3$	\leq	8

X^*	$= \{0, 2, 4\}$
SP^*	$= \{7.5, 0, 0.5\}$
OC^*	$= \{4.5, 0, 0\}$
SV^*	$= \{0, 1, 0\}$
Z^*	$= 34$

–Dual:





Min: $Y = 4W_1 + 7W_2 + 8W_3$			
s.t.	$W_1 + W_2$	\geq	3
	$W_2 + 2W_3$	\geq	1
	$W_1 + W_2 + W_3$	\geq	8

W^*	$= \{7.5, 0, 0.5\}$
dSV^*	$= \{4.5, 0, 0\}$
dSP^*	$= \{0, 2, 4\}$
dOC^*	$= \{0, 1, 0\}$
Y^*	$= 34$

Dual/Primal Solution Relationships

Primal

Dual

Decision Variables		Shadow Prices
Shadow Prices		Decision Variables
Opportunity Costs		Slack Variables
Slack Variables		Opportunity Costs