

## Linear Programming in Practice

▪ **Essential Issue:** To model non-linear reality with linear equations

- Activities
- Piece-wise linear approximations
- Fixed charges

▪ **Another practical question**

- Duality

## Activities

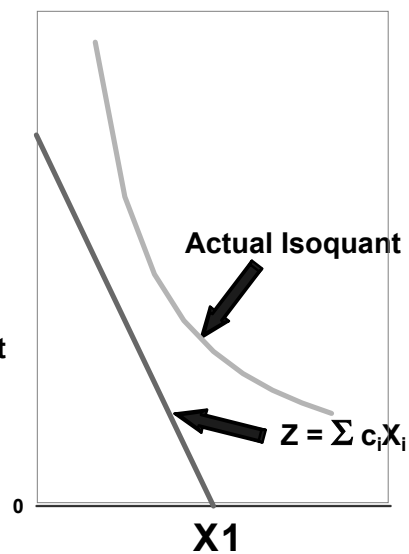
▪ **Motivation:**

▪ **If we use a standard production function**

$$f(\underline{X}) = \sum c_i X_i = Z$$

resources,  $\Rightarrow$  output

▪ **We are not able to represent typical production function with diminishing marginal returns and non-linear isoquants**



## Activities (continued)

### ▪Concept

- An activity is a
  - Specific way to use resources
  - in fixed proportions

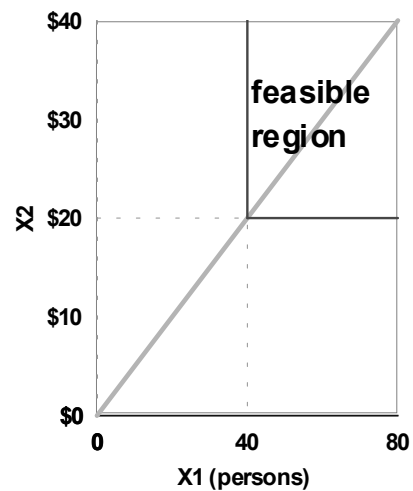
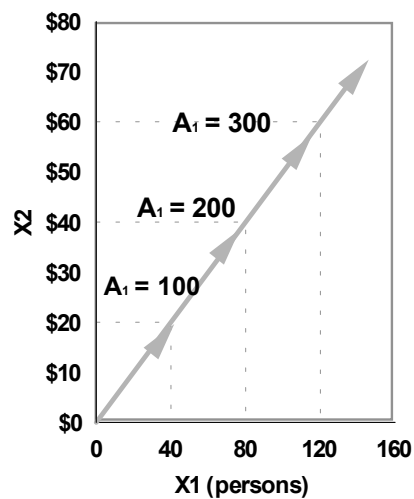
### ▪Physical interpretation is direct, e.g.:

- an aircraft using pilots, fuel / ton-km
- a machine requiring labor, materials per unit product

### ▪Think of activities as intermediates between resources and output

resources  $\rightarrow$  activities  $\rightarrow$  output

## Example: transport process $A_1$ uses 40 persons, \$20k to produce 100 T-km



## Activities

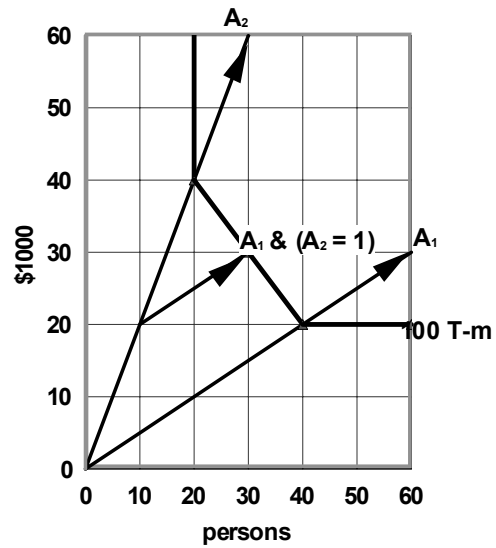
▪ Two Activities

▪  $A_1 = (40, 20K)$   
 $\implies 100 \text{ T-m}$

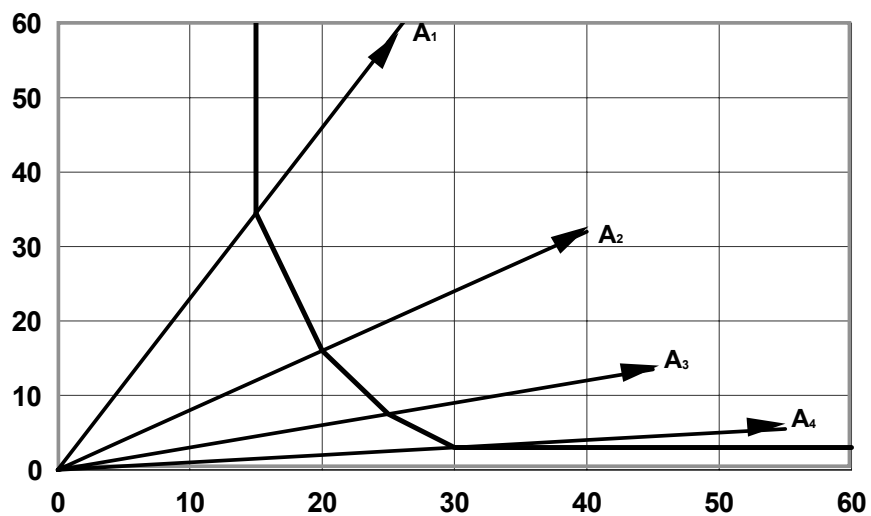
▪  $A_2 = (10, 20K)$   
 $\implies 50 \text{ T-m}$

$A_2 = 1 \{10, 20\} \implies 50$

$A_1 = \frac{1}{2} \{20, 10\} \implies 50$



## Many Activities



## LP Formulation with Activities

	Cr (kg)	C (kg)	Profit (\$)
<b>Process 1</b>	<b>6</b>	<b>4</b>	<b>30</b>
<b>Process 2</b>	<b>5</b>	<b>2</b>	<b>28</b>
<b>Process 3</b>	<b>3</b>	<b>6</b>	<b>29</b>

▪Optimize:  $Z = \sum c_i A_i$  -- subject to constraints

▪Example: Alloy optimization

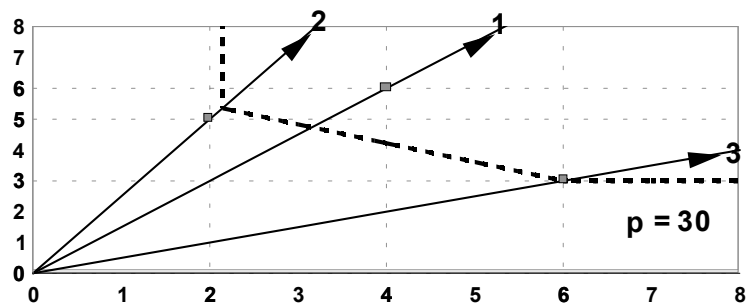
–Three possible processes, each with different unit costs

–subject to limits on Cr, C content

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## Formulation



$$\max Z = 30 P_1 + 28 P_2 + 29 P_3$$

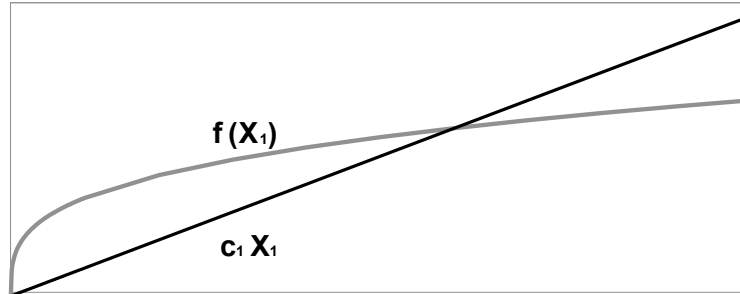
$$\text{s.t. } 6 P_1 + 5 P_2 + 3 P_3 \leq 26 \text{ (Cr)}$$

$$4 P_1 + 2 P_2 + 6 P_3 \leq 7 \text{ (C)}$$

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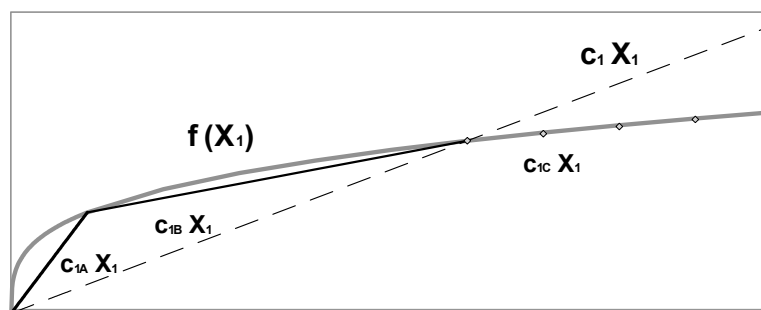
## Piece-Wise Linear Approximations (1)



### ▪Motivation:

- Returns to scale generally non-linear
- Straight line approximations are inaccurate

## Piece-Wise Linear Approximations (2)



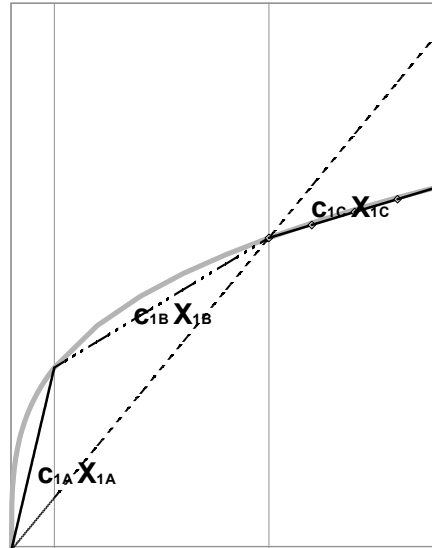
### ▪Concept:

- Represent  $f(X_1)$  with several lines

## Piece-wise Linear Approximations

### Implementation Notes:

- $X_1$  must be redefined as several variables -  $X_{1A}, X_{1B}, \dots$
- These new variables must not overlap, so  $X_{1A} < X_{1B}$ , etc.
- New variables and constraints make the LP larger and, therefore, more expensive



## Piece-wise Linear Approximations

### Mathematically:

▪ Given: 
$$\begin{aligned} \text{Max } Z &= f(X_1) + 4X_2 \\ \text{s.t.} \quad & 3X_1 + 6X_2 \leq 8 \end{aligned}$$

### Piece-wise linear approximation gives:

- $X_1 \Rightarrow X_{1A} + X_{1B}$
- $X_{1A}, X_{1B}$  have same  $a_{ij}$  as  $X_1$
- $c_1 = c_{1A}, c_{1B}$
- $X_{1A} < \text{cutoff } X \text{ value between } X_{1A} \text{ and } X_{1B}, X'$

▪ Thus: 
$$\begin{aligned} \text{Max } Z &= c_{1A} X_{1A} + c_{1B} X_{1B} + 4X_2 \\ \text{s.t.} \quad & 3 X_{1A} + 3X_{1B} + 6X_2 \leq 8 \\ & X_{1A} \leq X' \end{aligned}$$

## Piece-wise Linear Approximations (4)

- **Key Limitation:**

- ONLY works for convex feasible region!

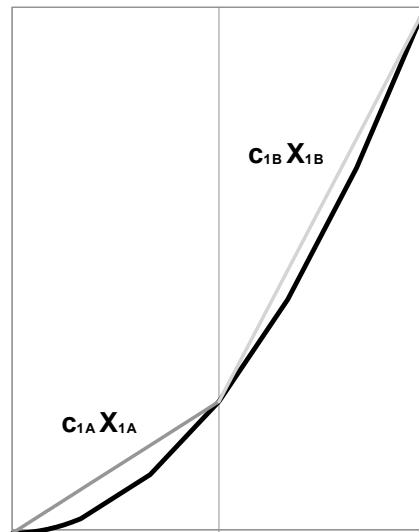
- **Why?**

- What if  $c_{1B} > c_{1A}$ ? (see fig)

$$\text{Max } Z = c_{1A} X_{1A} + c_{1B} X_{1B} + 4X_2$$

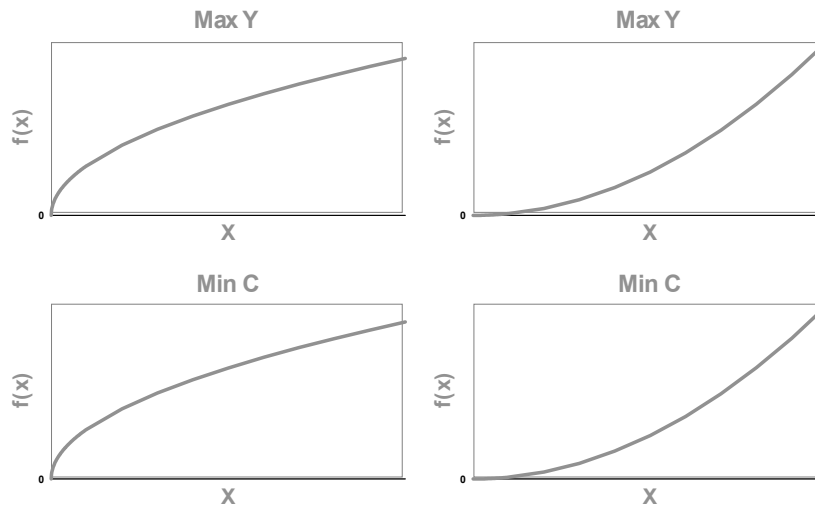
- **The LP will select  $X_{1B}$  before  $X_{1A}$ ;**

**Result may be meaningless!**



## Convex Feasible Regions Review:

Piecewise linear approximation works when FR is convex



## Fixed Charges

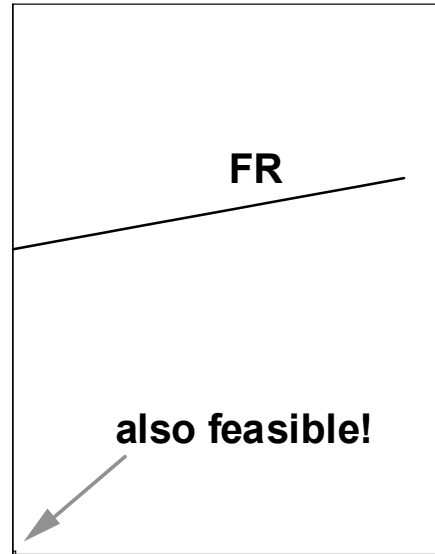
- Example: Warehousing
  - Cost = fixed rent, etc. + variable
  - Unless you choose not to operate it!

$$\begin{aligned} f(X_1) &= c_0 + c_1 X_1 & X_1 &\geq 0 \\ f(X_1) &= 0 & X_1 &= 0 \end{aligned}$$

- LP generally cannot handle fixed charges

### Exception:

- All  $X_i > 0$ ;  $X_i \neq 0$
- then subtract  $\sum c_0$  and optimize



## Duality

### ▪ Concept:

- A "dual" is a mirror-image form to another problem (the "primal")
- If primal = max; then dual = min
- If primal = min; then dual = max
- Dual contains all information of the primal, but in a different format
- Optimum value of primal = optimum value of dual

### ▪ Example:

- Primal: maximize output subject to budget limitations
- Dual: minimize costs subject to output requirements



## LP Duality

### ▪Mathematics:

–Given a Primal:

$$\begin{array}{ll} \text{Optimize:} & Z = \underline{c} X \\ \text{subject to:} & \underline{A} X \leq \geq \underline{B} \end{array}$$

–Dual is:

$$\begin{array}{ll} \text{Optimize:} & Y = \underline{B}^T W \\ \text{subject to:} & \underline{A}^T W \leq \geq \underline{c}^T \end{array}$$

### ▪Change of dimensionality between primal & dual:

– $\underline{c}^T$  and  $\underline{B}$  have different number of variables

### ▪Can use duality to:

- Reduce size of constraint matrix
- Speed up LP solution

## LP Duality - Example

$$\begin{array}{ll} \text{–Primal:} & \text{Max: } Z = X_1 + 2X_2 + 3X_3 \\ & \text{s.t. } 4X_1 + 2X_2 \leq 5 \\ & \quad \quad 6X_1 + 7X_2 + 9X_3 \leq 12 \end{array}$$

$$\begin{array}{ll} \text{–A =} & \begin{array}{ccc} 4 & 2 & 0 \\ 6 & 7 & 9 \end{array} & \text{A}^T = \begin{array}{c} 4 \ 6 \\ 2 \ 7 \\ 0 \ 9 \end{array} \end{array}$$

$$\begin{array}{ll} \text{–B =} & \begin{array}{c} 5 \\ 12 \end{array} & \text{B}^T = \begin{array}{cc} 5 & 12 \end{array} \end{array}$$

$$\begin{array}{ll} \text{–C =} & \begin{array}{ccc} 1 & 2 & 3 \end{array} & \text{C}^T = \begin{array}{c} 1 \\ 2 \\ 3 \end{array} \end{array}$$

## LP Duality - Example - continued

–Primal:            Max:  $Z = X_1 + 2X_2 + 3X_3$   
                          s.t.  $4X_1 + 2X_2 \leq 5$   
                                           $6X_1 + 7X_2 + 9X_3 \leq 12$

–Dual:            Min:  $Y = 5W_1 + 12W_2$   
                          s.t.  $4W_1 + 6W_2 \geq 1$   
                                           $2W_1 + 7W_2 \geq 2$   
                                                   $9W_2 \geq 3$

## LP Duality - Interpretation of Results

–Primal:

Max: $Z = 3X_1 + X_2 + 8X_3$			
s.t. $X_1 + X_3$	$\leq$	4	$X^* = \{0,2,4\}$ $SP^* = \{7.5,0,0.5\}$ $OC^* = \{4.5,0,0\}$ $SV^* = \{0,1,0\}$ $Z^* = 34$
$X_1 + X_2 + X_3$	$\leq$	7	
$2X_2 + X_3$	$\leq$	8	





–Dual:

Min: $Y = 4W_1 + 7W_2 + 8W_3$			
s.t. $W_1 + W_2$	$\geq$	3	$W^* = \{7.5,0,0.5\}$ $dSV^* = \{4.5,0,0\}$ $dSP^* = \{0,2,4\}$ $dOC^* = \{0,1,0\}$ $Y^* = 34$
$W_2 + 2W_3$	$\geq$	1	
$W_1 + W_2 + W_3$	$\geq$	8	

## Dual/Primal Solution Relationships

### Primal

### Dual

Decision Variables		Shadow Prices
Shadow Prices		Decision Variables
Opportunity Costs		Slack Variables
Slack Variables		Opportunity Costs