

## Constrained Optimization

### Outline

### Unconstrained Optimization (Review)

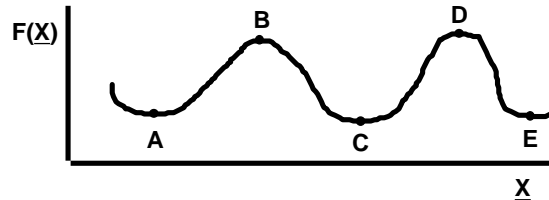
### Constrained Optimization

1. Approach
2. Equality constraints
  - Lagrangeans
  - Shadow prices
3. Inequality constraints
  - Kuhn-Tucker conditions
  - Complementary slackness

## Unconstrained Optimization (1)

- **Definitions:**
  - **Optimization** = Maximum of desired quantity  
= Minimum of undesired quantity
  - **Objective Function** = Expression to be optimized  
=  $Z(\underline{X})$
  - **Decision Variables** = Variables about which we can make decisions  
=  $\underline{X}$

## Unconstrained Optimization (2)



- By calculus:

If  $F(X)$  continuous, analytic:  
Condition for maxima and minima

$$\frac{\partial F(X)}{\partial X_i} = 0 \quad \forall_i$$

## Unconstrained Optimization (2) (cont'd)

- Secondary conditions:

$$\frac{\partial^2 F(X)}{\partial X_i^2} < 0 \quad \Rightarrow \text{Max} \quad (B,D)$$

$$\frac{\partial^2 F(X)}{\partial X_i^2} > 0 \quad \Rightarrow \text{Min} \quad (A,C,E)$$

## Unconstrained Optimization (3)

- Example: Housing insulation

$$F(x) = K_1 / x + K_2 x$$

Cost = Fuel cost + Insulation cost

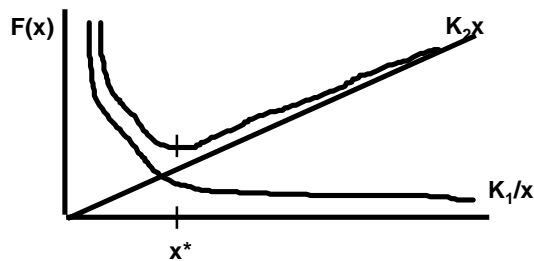
$x$  = Thickness of insulation

$$\partial F(x) / \partial x = 0 = -K_1 / x^2 + K_2$$

$$\Rightarrow x^* = \{K_1 / K_2\}^{1/2}$$

(starred quantities are optimal)

## Unconstrained Optimization (3) (cont'd)



## Constrained Optimization

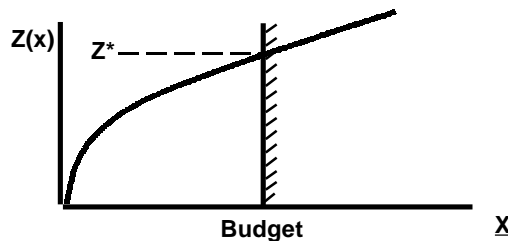
- Approach

To solve situations of increasing complexity (equality, inequality constraints) ...

Transform more difficult situation into one we know how to deal with

## Equality Constraints

- Example: Best use of budget
- Maximize: Output =  $Z(\underline{X}) = a_0 x_1^{a1} x_2^{a2}$
- Subject to (s.t.):  
Total costs = Budget =  $p_1 x_1 + p_2 x_2$



Note:  $\partial Z(\underline{X}) / \partial \underline{X} \neq 0$  at optimum

## Lagrangean Method (1)

Transforms equality constraints into unconstrained problem

Start with:

$$\text{Opt: } F(\mathbf{x})$$

$$\text{s.t.: } g_j(\mathbf{x}) = b_j \Rightarrow g_j(\mathbf{x}) - b_j = 0$$

Get to:

$$L = F(\mathbf{x}) - \sum_j \lambda_j [g_j(\mathbf{x}) - b_j]$$

$\lambda_j$  = Lagrangean multipliers  
(these are unknown quantities which must be solved for)

Note:  $[g_j(\mathbf{x}) - b_j] = 0$  so that  $\text{opt } F(\mathbf{x}) = \text{opt } L$

## Lagrangean Method (2)

- To optimize L:

$$\frac{\partial L}{\partial x_i} = 0 \quad \forall_i$$

$$\frac{\partial L}{\partial \lambda_j} = 0 \quad \forall_j$$

- Example:

$$\text{Opt: } F(\mathbf{x}) = 6x_1x_2$$

$$\text{s.t.: } g(\mathbf{x}) = 3x_1 + 4x_2 = 18$$

$$L = 6x_1x_2 - \lambda(3x_1 + 4x_2 - 18)$$

$$\frac{\partial L}{\partial x_1} = 6x_2 - 3\lambda = 0$$

$$\frac{\partial L}{\partial x_2} = 6x_1 - 4\lambda = 0$$

$$\frac{\partial L}{\partial \lambda} = 3x_1 + 4x_2 - 18 = 0$$

## Lagrangean Method (3)

- so that:  $\lambda = 2x_2 = 1.5x_1$   
 $x_2 = 0.75x_1$   
 $3x_1 + 3x_1 - 18 = 0 \Rightarrow x_1^* = 3$   
 $x_2^* = 2.25 \quad \lambda^* = 4.5$   
 $F(x)^* = 40.5$

## Shadow Prices

- **Definition:**
  - **Shadow Price = Rate of change of objective function per unit change of constraint**  
 $= \partial F(\underline{x}) / \partial b_j$
  - **This is meaning of Lagrangean multiplier**  
 $SP_j = \partial F(\underline{x})^* / \partial b_j = \lambda_j$
  - **Naturally, this is an instantaneous rate**

## Shadow Prices (cont'd)

- Example:

$$\text{Opt: } F(\mathbf{x}) = 6x_1x_2$$

$$\text{s.t.: } g(\mathbf{x}) = 3x_1 + 4x_2 = 18.1$$

$$\Rightarrow x_1^* = 18.1/6 \qquad x_2^* = 18.1/8$$

$$F(\mathbf{x})^* = 6(18.1/6)(18.1/8) = 40.95$$

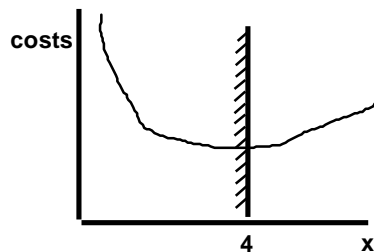
$$\Delta F(\mathbf{x}) = 40.95 - 40.5 = 0.45 = \hat{\lambda}^* (0.1)$$

## Inequality Constraints

- Example: Housing insulation

$$\text{Min: Costs} = K_1 / x + K_2x$$

$$\text{s.t.: } x \geq 4 \text{ (minimum thickness)}$$



## Inequality Constraints (1)

- Kuhn-Tucker conditions
- Transform inequalities into equalities, then proceed as before
- Again, introduce new variable

## Inequality Constraints (2)

- This is “slack variable”,  $s_j$   
 $g_j(\underline{x}) \leq b_j \Rightarrow g_j(\underline{x}) + s_j^2 = b_j$   
 $g_j(\underline{x}) \geq b_j \Rightarrow g_j(\underline{x}) + s_j^2 = b_j$

start from:

opt:  $F(\underline{x})$

s.t.:  $g_j(\underline{x}) \leq b_j$

get to:

$$L = F(\underline{x}) - \sum_j \lambda_j [g_j(\underline{x}) + s_j^2 - b_j]$$



## Inequality Constraints (3)

- Complementary slackness
- The optimality conditions are:

$$\partial L / \partial x_i = 0$$

$$\partial L / \partial \lambda_j = 0$$

$$\text{plus: } \partial L / \partial s_j = 2s_j\lambda_j = 0$$

- These new equations imply:

$$s_j = 0 \qquad \lambda_j \neq 0$$

or

$$s_j \neq 0 \qquad \lambda_j = 0$$

**This is the “complementary slackness”**

## Inequality Constraints (4)

- Interpretation:
- If there is slack on  $b_j$ ,  
(i.e. more than enough of it)  
=> No value to objective function  
to having more:  $\lambda_j = \partial F(\underline{x}) / \partial b_j = 0$

If  $\lambda_j \neq 0$ , then all available  $b_j$  used  
=>  $s_j = 0$